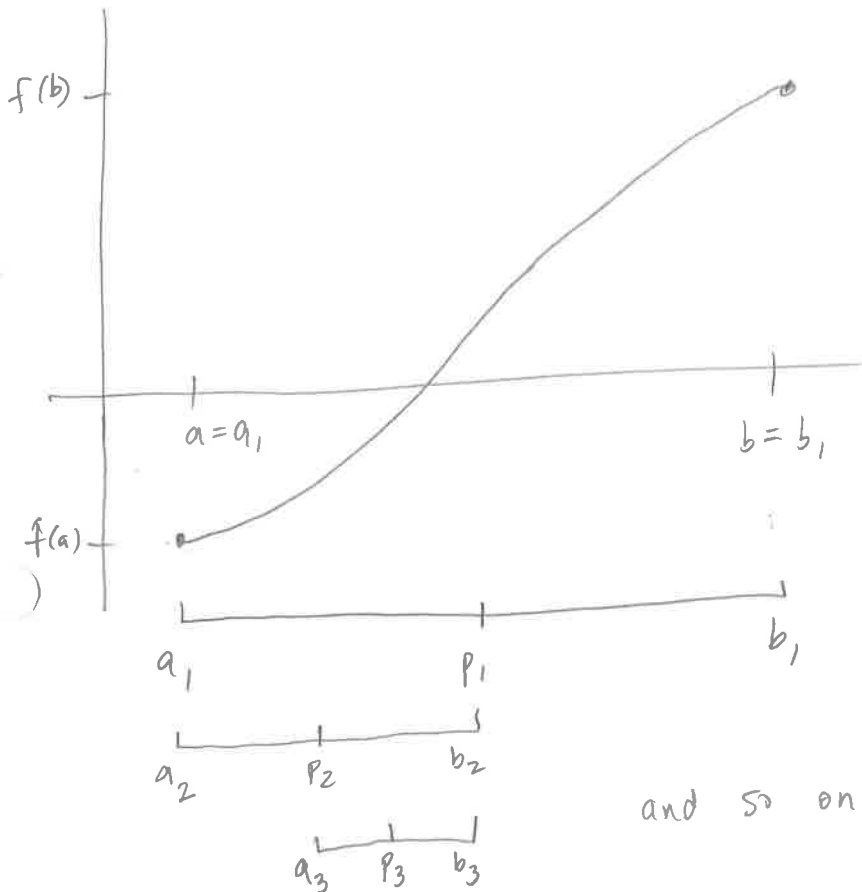


## 2.1 Bisection Method (Binary Search)

Play High/Low.



### Algorithm

Input  $a, b, \text{TOL}, \text{max \# of iterations } N_0$

Output approx solution  $p$  or message of failure

Step 1 for  $(i=1 \text{ to } N_0)$  do step 2-4.

Step 2 set  $p = \frac{a+b}{2}$

Step 3 if  $f(p)=0$  or  $\frac{(b-a)}{2} < \text{TOL}$  then  
output( $p$ );  
Stop

Step 4 If  $f(a)f(p) > 0$  then set  $a=p$

else set  $b=p$

Step 5 Output ('Method failed after  $N_0$  iteration!') Stop.

$$\text{Try on } f(x) = x^2 - 2x + 1 = 0.$$

### Thm 2.1

Let  $f \in C[a, b]$  and suppose  $f(a) \cdot f(b) < 0$ .

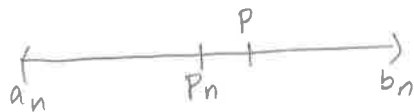
The Bisection method (Alg. 2.1) generates a sequence  $\{P_n\}$  approximating  $p$  with the property

$$|P_n - p| \leq \frac{b-a}{2^n}, \quad n \geq 1$$

Proof:  $\forall n \geq 1$ , we have

$$b_n - a_n = \frac{1}{2^{n-1}}(b-a) \quad (\text{We half the interval each step.})$$

Note that  $p \in (a_n, b_n)$ .



Since  $P_n = \frac{1}{2}(a_n + b_n)$ ,  $\forall n \geq 1$ , it follows that

$$|P_n - p| \leq \frac{1}{2}(b_n - a_n) = \frac{1}{2} \frac{(b-a)}{2^{n-1}} = \frac{b-a}{2^n}$$

Thus, as  $n \rightarrow \infty$ ,  $P_n \rightarrow p$  at the rate of  $O(\frac{1}{2^n})$ .

Ex: How many iteration of Bisection would be required to get the approx accurate to within  $10^{-8}$   
Suppose  $a=0, b=1$ . Then

$$|P_n - p| \leq \frac{b-a}{2^n} < 10^{-8}$$

$$2^{-n} < 10^{-8}$$

$$-n \log 2 < \log 10^{-8}$$

$$n > \frac{8}{\log 2} = 26.57 \Rightarrow 27 \text{ iterations}$$